Efficient Tensor Contraction Algorithms for Coupled Cluster

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- Motivation
- Interface
- Coupled Cluster with CTF
- Internal mechanism
- Performance

Symmetry Preserving Algorithm

- Instances in matrix computations
- General symmetric contractions
- Application to coupled-cluster

Conclusion

Motivation

The problem

We want portable infrastructure and scalable algorithms for tensor-based electronic structure methods

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- a 'stable' solution must provide a high-level abstraction that permits rapid manipulation of the algebra
- scalability must be achieved both for intranode (shared memory) and internode (distributed memory) parallelism

Cyclops (cyclic operations) Tensor Framework (CTF)

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- optimized for distributed networks, shared memory, and accelerators
- open source, BSD license, https://github.com/solomonik/ctf

CTF is orchestrated by bulk synchronous operations on a set of processors

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- an 'AS' dimension is antisymmetric with the next (also 'SY' and 'SH')
- tensors are templated by the element type (double by default)
- custom algebraic structures (set, group, monoid, semiring, ring) may be defined by the user

CTF can express a tensor contraction like

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- Z, F, T, W should all be defined on the same world and all processes in the world must call the contraction bulk synchronously
- user-defined (mixed-type) scalar tensor functions can be applied instead of + and *

Quantum chemistry codes using CTF

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- Libtensor has been integrated with CTF by Evgeny Epifanovsky
- **Q-Chem** can leverage Libtensor and integration with CTF is almost complete

CCSD

Extracted from Aquarius (Devin Matthews' code, https://github.com/devinamatthews/aquarius)

FMI["mi"] += 0.5*WMNEF["mnef"]*T(2)["efin"]; WMNIJ["mnij"] += 0.5*WMNEF["mnef"]*T(2)["efij"]; FAE["ae"] -= 0.5*WMNEF["mnef"]*T(2)["afmn"]; WAMEI["amei"] -= 0.5*WMNEF["mnef"]*T(2)["afin"];

CCSDT

Extracted from Aquarius (Devin Matthews' code)

```
Z(1)["ai"] += 0.25*WMNEF["mnef"]*T(3)["aefimn"];
```

```
Z(2)["abij"] += 0.5*WAMEF["bmef"]*T(3)["aefijm"];
Z(2)["abij"] -= 0.5*WMNEJ["mnej"]*T(3)["abeinm"];
Z(2)["abij"] += FME["me"]*T(3)["abeijm"];
```

```
Z(3)["abcijk"] = WABEJ["bcek"]*T(2)["aeij"];
Z(3)["abcijk"] -= WAMIJ["bmjk"]*T(2)["acim"];
Z(3)["abcijk"] += FAE["ce"]*T(3)["abeijk"];
Z(3)["abcijk"] -= FMI["mk"]*T(3)["abcijm"];
Z(3)["abcijk"] += 0.5*WABEF["abef"]*T(3)["efcijk"];
Z(3)["abcijk"] += 0.5*WMNIJ["mnij"]*T(3)["abcmnk"];
Z(3)["abcijk"] -= WAMEI["amei"]*T(3)["ebcmjk"];
```

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- extract permuted tensor slices (e.g. arbitrary subsets of rows and columns)

CTF tensor decomposition

• cyclic layout used to preserve packed symmetric structure

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Tensor decomposition and mapping

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CTF mapping logic

- arrange physical topology into all possible processor grids
- for each contraction autotune over all topologies and mappings
- select best mapping based on performance models (communication cost, memory requirements, etc.)

Symmetric matrix representation

Symmetric matrix



Unique part of symmetric matrix



Internal mechanism

Blocked distributions of a symmetric matrix

Naive blocked layout

Block-cyclic layout



Cyclic distribution of a symmetric matrix

Cyclic layout ~ Improved blocked layout



Tensor contraction mapping visualization



The following three redistribution kernels are provided by CTF

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 - processors exchange blocks via point-to-point messages

Comparison with NWChem

NWChem is a commonly-used distributed-memory quantum chemistry method suite

- provides CCSD and CCSDT
- uses Global Arrays (GA) tensor partitioning and contraction
- Tensor Contraction Engine (TCE) factorizes CC equations and generated GA code



Coupled-cluster code on BlueGene/Q (Mira)

CCSD up to 55 water molecules with cc-pVDZ CCSDT up to 10 water molecules with cc-pVDZ



Coupled-cluster code on Cray XC30 (Edison)

CCSD up to 50 water molecules with cc-pVDZ CCSDT up to 10 water molecules with cc-pVDZ



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 $c_i = \sum_{j=1}^n W_{ij}$

which requires n^2 multiplications and n^2 additions

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• The *symmetry preserving algorithm* employs a *symmetric* intermediate matrix **Z**,

$$Z_{ij} = A_{ij} \cdot (b_i + b_j)$$
 $c_i = \sum_{j=1}^n Z_{ij} - \left(\sum_{j=1}^n A_{ij}\right) \cdot b_i$

which requires $\frac{n^2}{2}$ multiplications and $\frac{5n^2}{2}$ additions

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- $\mathbf{C} = \mathbf{A} \cdot \mathbf{B} + \mathbf{B} \cdot \mathbf{A}$ is usually computed via a nonsymmetric intermediate order 3 tensor \mathbf{W} ,

$$W_{ijk} = A_{ik} \cdot B_{kj}$$
 $\overline{W}_{ij} = \sum_{k} W_{ijk}$ $C_{ij} = W_{ij} + W_{ji}$.

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which requires n^3 multiplications and n^3 additions.

 The symmetry preserving algorithm employs a symmetric intermediate tensor Z using n³/6 multiplications and 7n³/6 additions,

$$Z_{ijk} = (A_{ij} + A_{ik} + A_{jk}) \cdot (B_{ij} + B_{ik} + B_{jk}) \qquad v_i = \sum_{k=1}^n A_{ik} \cdot B_{ik}$$
$$C_{ij} = \sum_{k=1}^n Z_{ijk} - n \cdot A_{ij} \cdot B_{ij} - v_i - v_j - \left(\sum_{k=1}^n A_{ik}\right) \cdot B_{ij} - A_{ij} \cdot \left(\sum_{k=1}^n B_{ik}\right)$$

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- Communication cost lower and upper bounds derived

Application to CCSD

The CCSD contraction

$$Z^{aar{k}}_{iar{c}} = \sum_{b}\sum_{j}T^{ab}_{ij}\cdot V^{jar{k}}_{bar{c}}$$

usually requires $2n^6$ total operations.

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The symmetry-preserving algorithm can be applied over the indices

$$\mathsf{Z}^{\mathsf{a}} = \sum_{b} \mathsf{T}^{\mathsf{a}b} \cdot \mathsf{V}_{\mathsf{b}}$$

with each multiplication being a contraction over the other four indices i, j, \bar{c}, \bar{k} , which is more expensive than the addition operations, yielding n^6 operations to leading order.

Application to CCSD(T) and CCSDT(Q)

The CCSD(T) contraction

$$T^{abar{c}}_{ijar{k}} = P(a,b)P(i,j)\sum_{ar{l}=1}^n T^{aar{c}}_{iar{l}}\cdot W^{ar{l}b}_{jar{k}}$$

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$$\mathbf{T}^{\mathbf{ab}} = P(\mathbf{a}, \mathbf{b})\mathbf{T}^{\mathbf{a}} \cdot \mathbf{W}^{\mathbf{b}}$$
 and $\mathbf{T}_{\mathbf{ij}} = P(i, j)\mathbf{T}_{\mathbf{i}} \cdot \mathbf{T}_{\mathbf{j}}$

with each multiplication in the latter being a contraction over the remaining three indices \bar{c}, \bar{k} , and \bar{l} , for a total of $n^7/2$ leading order operations.

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For a similar CCSDT(Q) contraction, which usually requires $n^9/2$ operations, the symmetry preserving algorithm achieves $n^9/18$.

Conclusion

Future work on symmetry-preserving algorithms

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- iterative performance-model refinement via online learning
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- sparse tensors

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